# CSCI 3110 Tutorial 5 

## Reviewed June 14, 2019

1. Consider the task of sorting an unsorted array $A[1 \ldots n]$ : a task we can perform by merge sort in time $\mathrm{O}(n \lg n)$. Show that any algorithm that accesses the array only via comparisons (that is, by asking questions of the form "is $A[i] \leq z$ ?"), must take $\Omega(n \lg n)$ steps.
2. Given a sequence of real numbers, $X=\left(x_{1}, x_{2}, \ldots, x_{n}\right)$, an exchanged pair in $X$ is a pair $\left(x_{i}, x_{j}\right)$ such that $i<j$ and $x_{i}>x_{j}$. Note that an element $x_{i}$ can be part of up to $n-1$ exchanged pairs. In particular, the maximal possible number of exchanged pairs in $X$ is $\frac{n(n-1)}{2}$, which is achieved if the array is sorted in descending order. Assume that the sequence $X$ is stored in an array of size $n$. Develop a divide-and-conquer algorithm that counts the exchanged pairs in $X$. (Your algorithm should take $\mathrm{O}(n \lg n)$ time.) Prove that your algorithm is correct. Argue briefly why your algorithm takes $\mathrm{O}(n \lg n)$ time.
