

CSCI 3110

Fun with Algorithms

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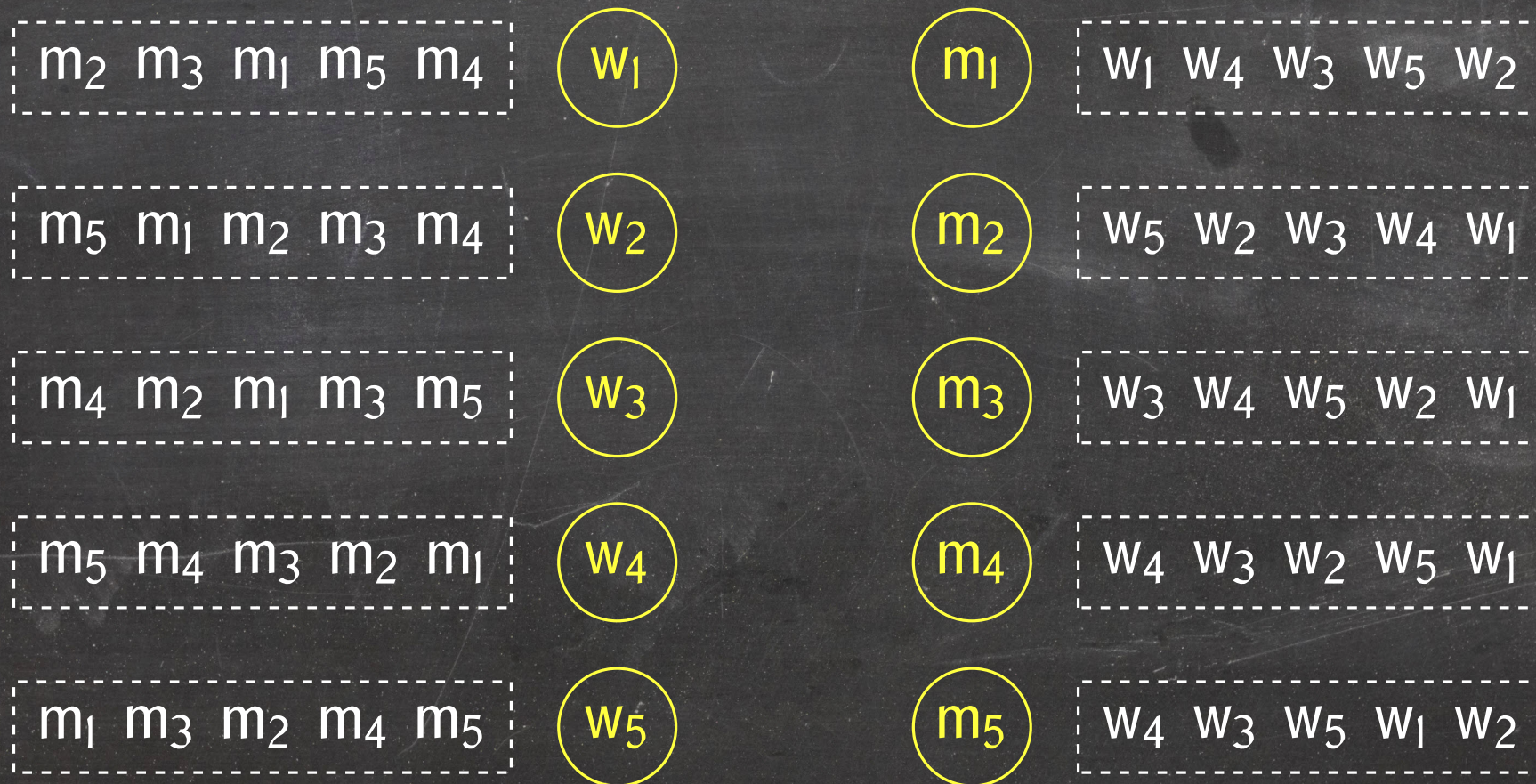
Faculty of Computer Science
Dalhousie University
Summer 2019

(Adapted from Slides by Norbert Zeh)

Stable Matching: An Introductory Example

Given:

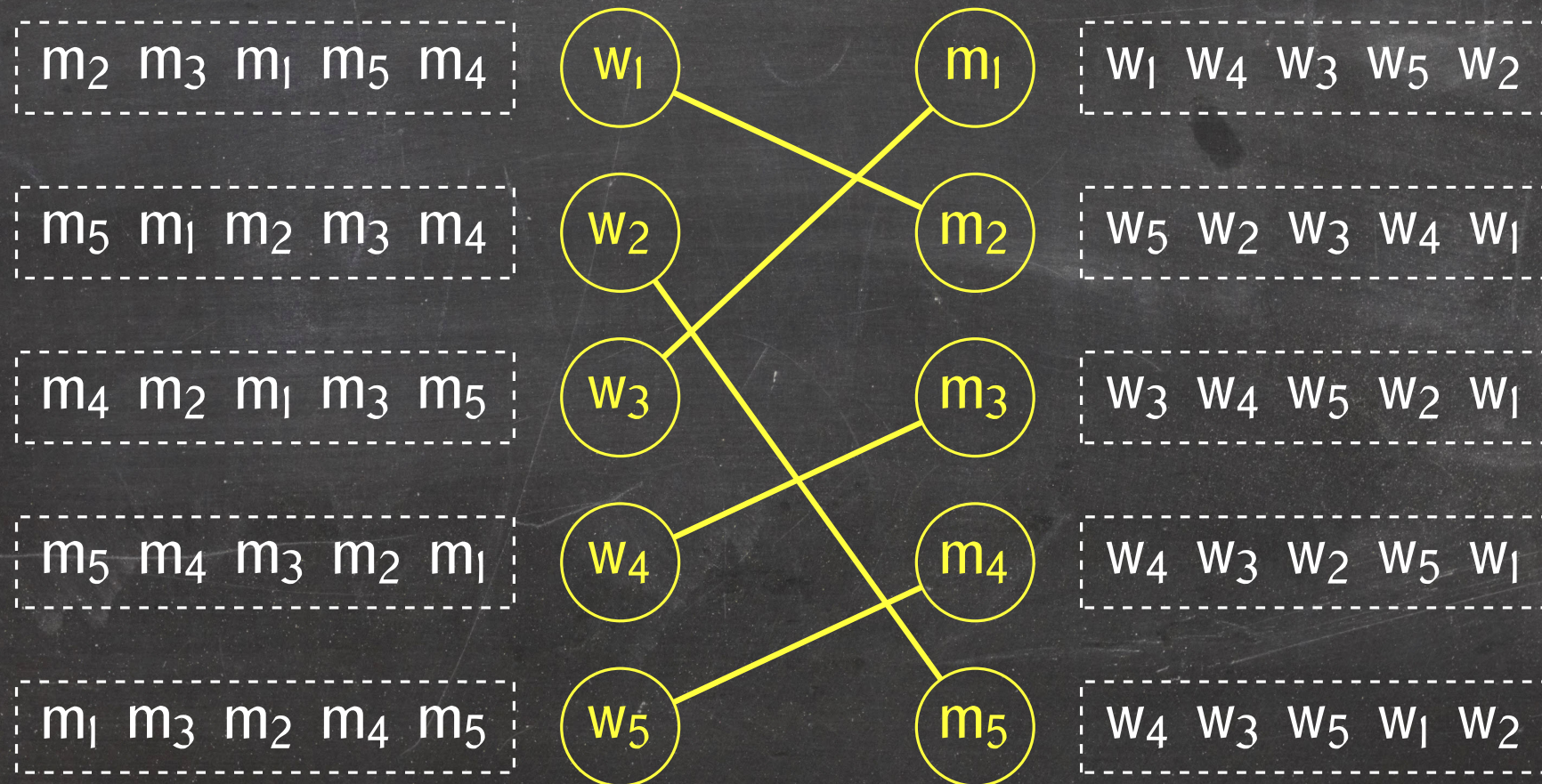
- n women w_1, w_2, \dots, w_n
- n men m_1, m_2, \dots, m_n
- A preference list for each



Stable Matching: An Introductory Example

Output:

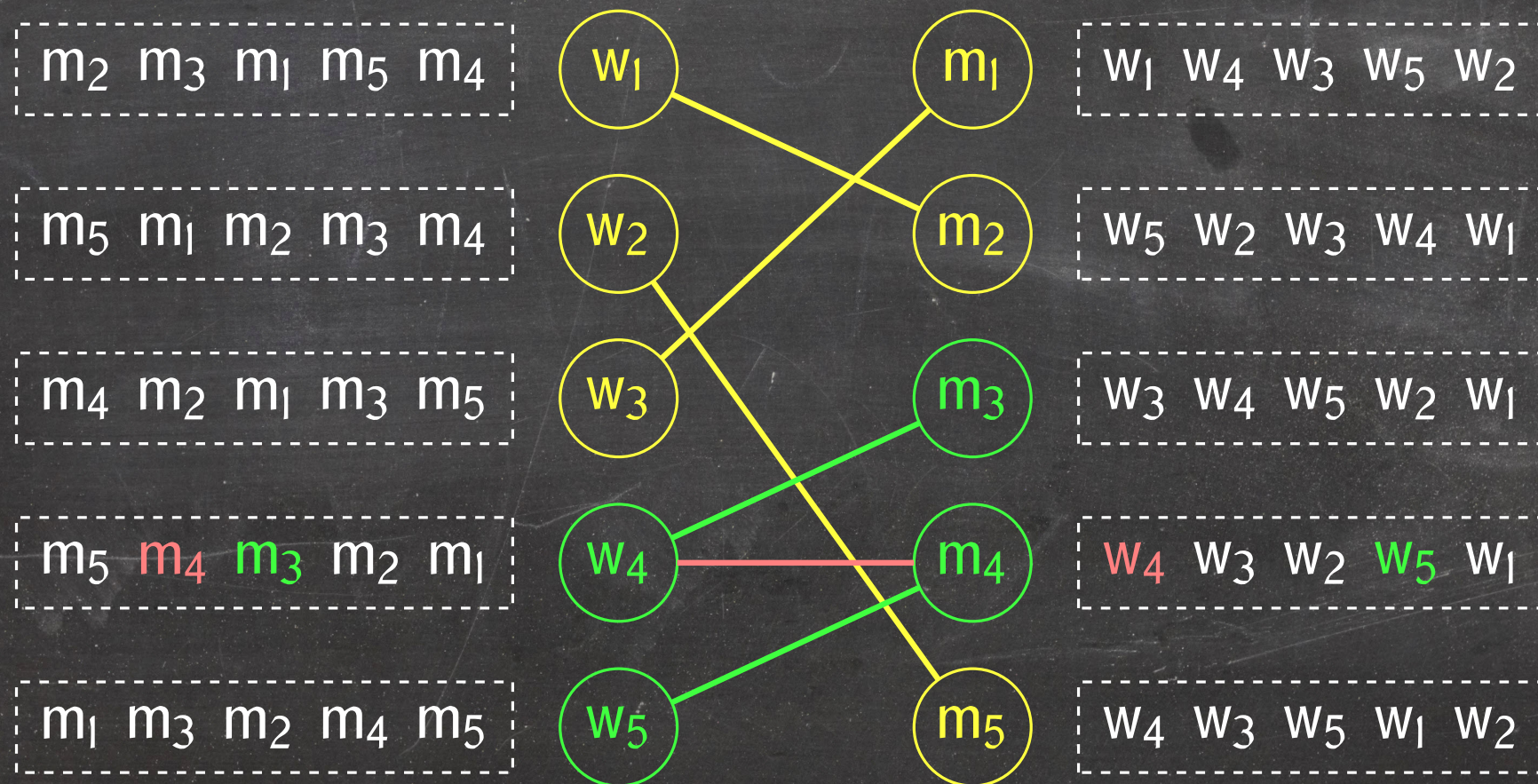
- A set of n marriages $\{(w_{i_1}, m_{j_1}), (w_{i_2}, m_{j_2}), \dots, (w_{i_n}, m_{j_n})\}$
- Every man is married
- Every woman is married
- The marriages are **stable**



Stable Matching: An Introductory Example

A pair of marriages (m, w) and (m', w') is **unstable** if

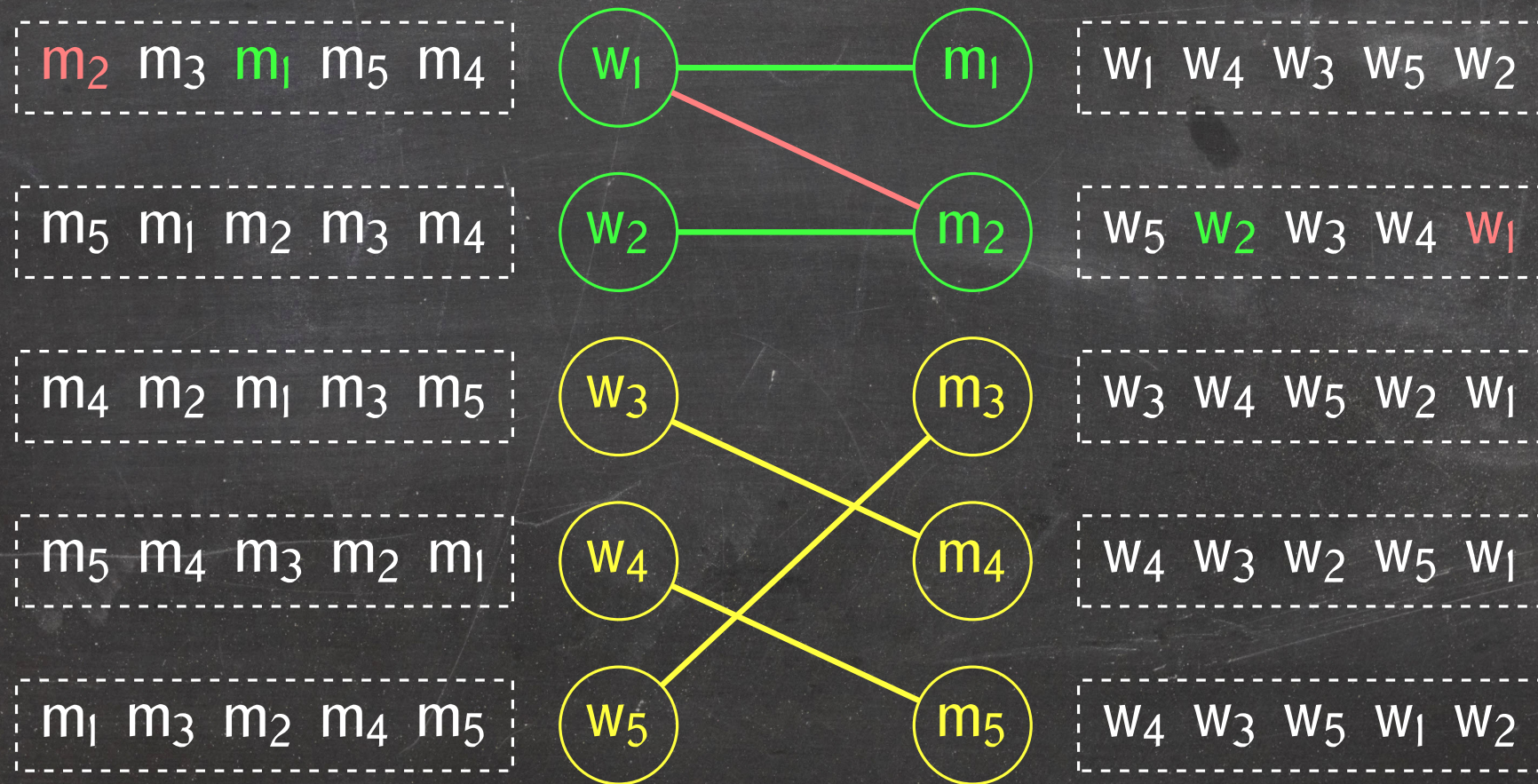
- w prefers m' over m ($m' \prec_w m$)
- m' prefers w over w' ($w \prec_{m'} w'$)



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Stable Matching: A Solution Inspired By Real Life

StableMatching(M, W)

```
1  while there exists an unmarried man m
2    do m proposes to the most preferable woman w he has not proposed to yet
3    if w is unmarried or likes m better than her current partner m'
4      then if w is married
5          then w divorces m'
6          w marries m
```


Stable Matching: A Solution Inspired By Real Life

StableMatching(M, W)

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Questions we can and should ask about the algorithm:

- Is there always a stable matching?
- Does the algorithm always terminate?
- Does the algorithm always produce a stable matching?
- How efficient is the algorithm? Can we bound its running time?

Course Outline

- Correctness proofs
- Analysis of resource consumption
- Algorithm design techniques
 - Graph exploration
 - Greedy algorithms
 - Divide and conquer
 - Dynamic programming
 - Data structuring
 - Randomization
- NP-completeness and intractability

General Information

Instructor: Christopher Whidden

Office: CS 315

Office hours: Thursday 1:00–2:00

Email: cwhidden@dal.ca

Textbook: Cormen, Leiserson, Rivest, Stein. **Introduction to Algorithms.**
3rd edition, MIT Press, 2009.

Zeh. **Data Structures.**

CSCI 3110 Lecture Notes, 2005.

Website: <http://www.cs.dal.ca/~whidden/CSCI3110>

TAs: Yuhan Fu

Mozhgan Saeid

Younan Gao

Midterm: July 4

Grading

- 10 Assignments (A)

The best 8 count. Each carries equal weight.

- Midterm (M)
- Final (F)

$$\text{Final grade} = \max \left(\begin{array}{l} 60\% \cdot F + 40\% \cdot A \\ 40\% \cdot F + 20\% \cdot M + 40\% \cdot A \end{array} \right)$$

Collaboration, Plagiarism, Late Assignments

Collaboration

- Groups of up to three people are allowed to collaborate on assignments.
- Every group hands in one set of solutions; every group member gets the same marks.
- Collaboration between groups is not allowed!

Plagiarism

- Plagiarism will not be tolerated.
- Collaboration between groups is a form of plagiarism.

Late assignments

... will not be accepted. Assignments missed for a reason documented by a Student Declaration of Absence will be covered by your final exam score.

Please see course website for a detailed discussion of these rules.

Things I Expect You To Know

- Basic rules concerning logarithms
- Basic rules concerning limits
- Basic derivatives
- Propositional logic
- Elementary combinatorics (counting permutations, combinations, ...)
- Elementary probability theory (linearity of expectation, ...)
- Elementary data structures (arrays, lists, stacks, queues, ...)
- Standard sorting algorithms (insertion sort, quick sort, merge sort)
- Binary heaps