

# Midterm Review

Tuesday, July 2, 2019 11:38 AM

Fundamentals  
asymptotic notation

$o()$   $O()$   $\Omega$   $\Theta$

limits to prove  $f(n) = o(g(n))$

provide constants to prove  $f(n) = O(g(n))$   
 $\Omega$   
 $\Theta$

## Algorithm Analysis

What does a specific algorithm do?

analyze the running time of an algorithm

- Counting loops and operations
- bounding the number of function calls

BFS

- recurrence relations

recurrence relation

$$T(n) = 2T(n/2) + n \quad \text{mergesort}$$

master theorem

substitution

recursion tree

• Amortized analysis

Prim's algorithm  $O(m + n \log n)$

amortized using thin heap  
potential functions

prove correctness

termination

• usually simple

• bounded number of steps

e.g. looking at every element of array

look at any one element at most  $\lg n$  times

Correctness

• contradiction

- contradiction  
assume simplest case
- induction  
base case  
inductive step  
termination
- Stay ahead arguments (greedy)  
Similar to induction  
on termination  
argue your set size is the same  
length as optimal
- loop invariants  
form of induction  
show that an "invariant" holds  
at every step of a loop

# Greedy Algorithms ...

# Greedy Algorithms

make progress with local choices to obtain global optimum

problems: interval scheduling

MST - Kruskal

union find

Huffman coding

priority

priority queues (e.g.  $\text{Heap}$ )

Dijkstra's algorithm

techniques  
induction  
stay-ahead

# Graph Exploration

definitions

vertices, edges, adjacency list

undirected / directed

proofs: contradiction

BFS / DFS as building blocks

connected components

bipartiteness

top sort

strongly connected components

# Divide and Conquer

recursively break problem into smaller

subproblems, solve, and combine the solutions

e.g. mergesort, quicksort

techniques: induction

techniques: induction

recurrence relations

$$T(n) = aT\left(\frac{n}{b}\right) + f(n)$$

↑  
running time

↑  
#subproblems

↑  
size of subproblems

↑  
local work

mergesort  $T(n) = 2T\left(\frac{n}{2}\right) + \Theta(n)$

↓

closed form  $T(n) = \Theta(n \log n)$

problems: selection, matrix multiplication, closest pair

## Dynamic Programming

recursively break into smaller subproblems

techniques: recurrence relations

↓ provide answer, not running time

problems

chain matrix multiplication  
weighted interval scheduling  
sequence alignment  
shortest paths

shortest paths

how to handle negative weights

Floyd-Warshall,

all-pairs shortest paths