

An Augmenting Paths Algorithm That Terminates

CSCI 4113/6101



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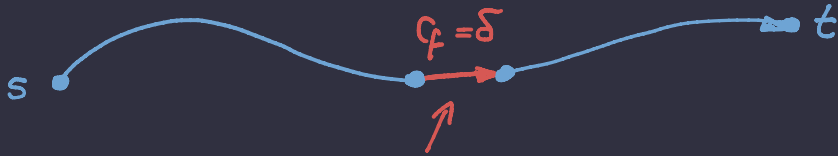
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- Uses BFS to find a shortest augmenting path in each iteration.

Lemma: Edmonds-Karp terminates after $O(nm)$ iterations.

Corollary: Edmonds-Karp finds a maximum st-flow in $O(nm^2)$ time.

Critical Edges

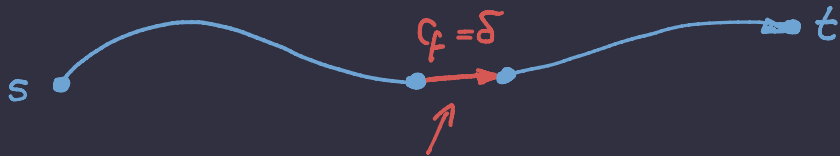
Augmenting path in G_f



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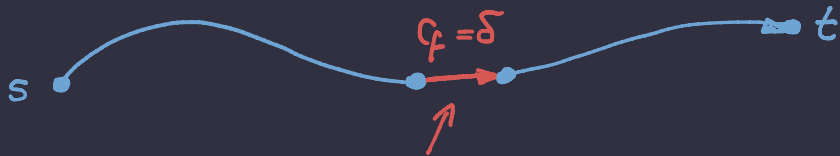


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Claim: Every edge is critical for at most $\frac{n}{2}$ augmenting paths.

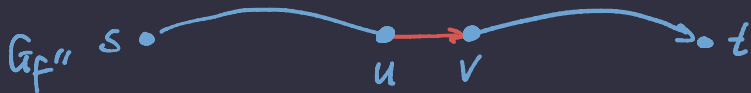
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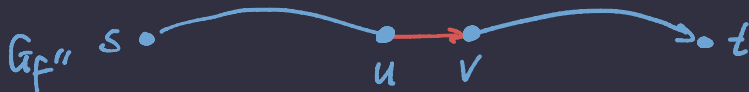
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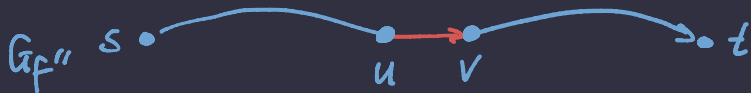


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$$\begin{aligned} d_{G_f'}(s, u) &= d_{G_f'}(s, v) + 1 \\ &\geq d_{G_f}(s, v) + 1 \end{aligned}$$



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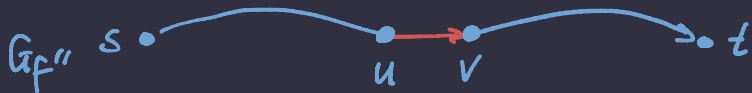
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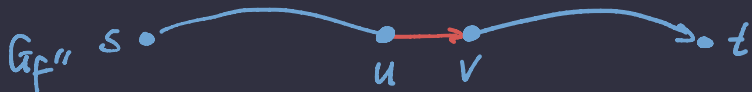
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$\Rightarrow (u, v)$ critical for $\leq \frac{n}{2}$ paths

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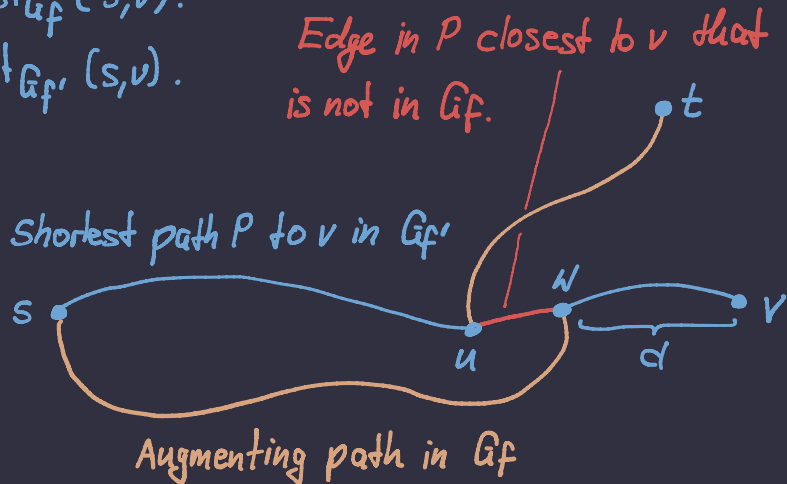
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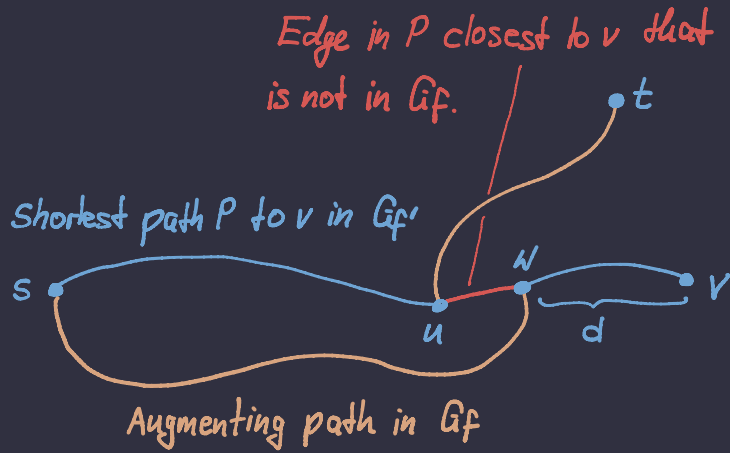
Choose v with minimum $\text{dist}_{G_{f'}}(s, v)$.

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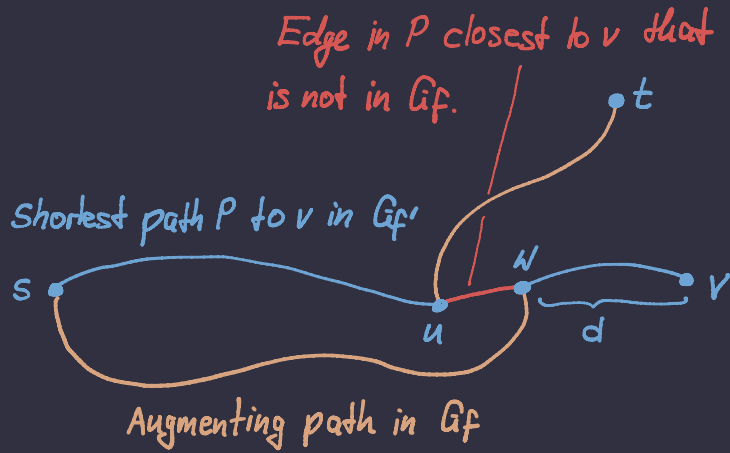
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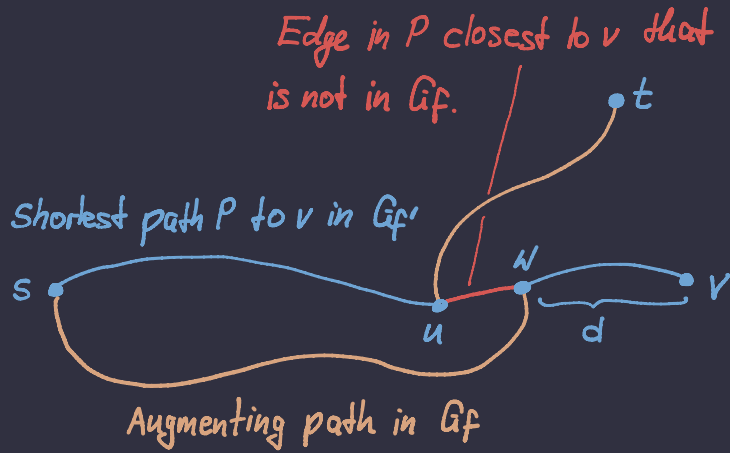




$$\begin{aligned} \text{dist}_{G_f'}(s, v) &= \text{dist}_{G_f'}(s, u) + d + 1 \\ &\geq \text{dist}_{G_f}(s, u) + d + 1 \end{aligned}$$



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$$\Rightarrow \text{dist}_{G_f}(s, v) < \text{dist}_{G_f'}(s, v)$$