

# Artificial Intelligence: Search

## Part 3: Objective optimization

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Based on the slides provided by Russell and Norvig, Chapter 4, Section 3–4

# Outline

- ◇ Hill-climbing
- ◇ Simulated annealing
- ◇ More formally on local search in continuous spaces

# Iterative improvement algorithms

In many optimization problems, **path** is irrelevant;  
the goal state itself is the solution

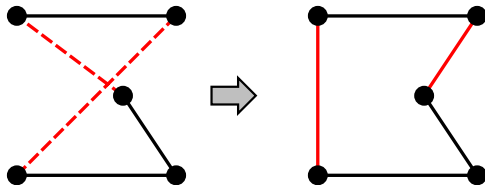
Then state space = set of “complete” configurations;  
find **optimal** configuration, e.g., TSP  
or, find configuration satisfying constraints, e.g., timetable

In such cases, can use **iterative improvement** algorithms;  
keep a single “current” state, try to improve it

Constant space, suitable for online as well as offline search

# Example: Travelling Salesperson Problem

Start with any complete tour, perform pairwise exchanges

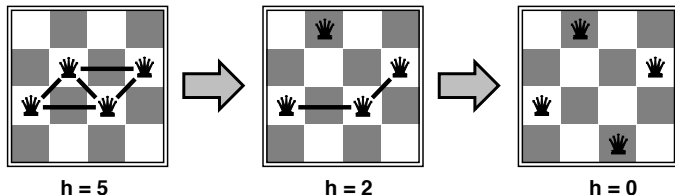


Variants of this approach get within 1% of optimal very quickly with thousands of cities

## Example: $n$ -queens

Put  $n$  queens on an  $n \times n$  board with no two queens on the same row, column, or diagonal

Move a queen to reduce number of conflicts



Almost always solves  $n$ -queens problems almost instantaneously for very large  $n$ , e.g.,  $n = 1\text{ million}$

# Hill-climbing (or gradient ascent/descent)

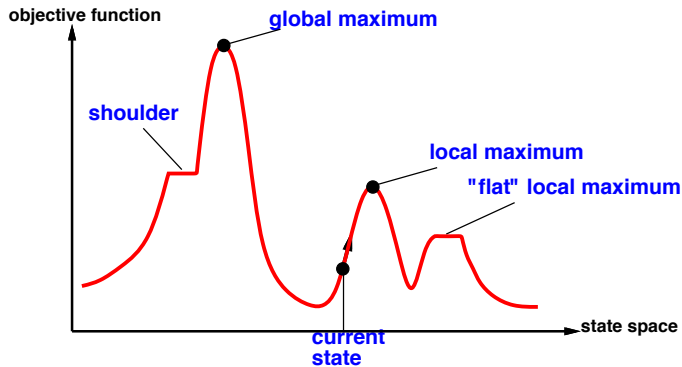
“Like climbing Everest in thick fog with amnesia”

```
function HILL-CLIMBING(problem) returns a state that is a local maximum
inputs: problem, a problem
local variables: current, a node
                   neighbor, a node

current ← MAKE-NODE(INITIAL-STATE[problem])
loop do
  neighbor ← a highest-valued successor of current
  if VALUE[neighbor] ≤ VALUE[current] then return STATE[current]
  current ← neighbor
end
```

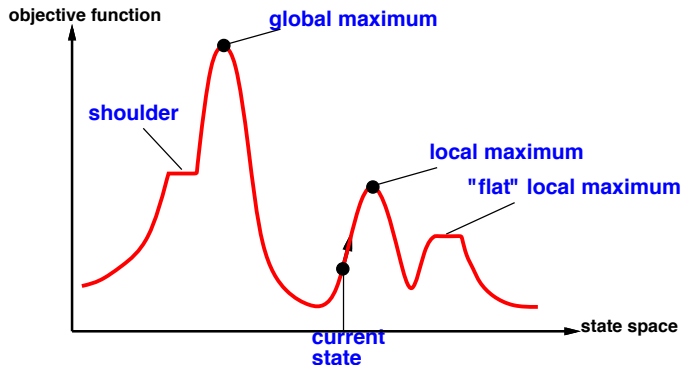
# Hill-climbing contd.

Useful to consider **state space landscape**



# Hill-climbing contd.

Useful to consider **state space landscape**



**Random-restart hill climbing** overcomes local maxima—trivially complete

**Random sideways moves** 🤪 escape from shoulders 🌀 loop on flat maxima



# Simulated annealing

Idea: escape local maxima by allowing some “bad” moves  
**but gradually decrease their size and frequency**

**function** SIMULATED-ANNEALING(*problem*, *schedule*) **returns** a solution state

**inputs:** *problem*, a problem

*schedule*, a mapping from time to “temperature”

**local variables:** *current*, a node

*next*, a node

*T*, a “temperature” controlling prob. of downward steps

*current* ← MAKE-NODE(INITIAL-STATE[*problem*])

**for**  $t \leftarrow 1$  **to**  $\infty$  **do**

*T* ← *schedule*[*t*]

**if** *T* = 0 **then return** *current*

*next* ← a randomly selected successor of *current*

$\Delta E \leftarrow \text{VALUE}[\textit{next}] - \text{VALUE}[\textit{current}]$

**if**  $\Delta E > 0$  **then** *current* ← *next*

**else** *current* ← *next* only with probability  $e^{\Delta E/T}$

# Properties of simulated annealing

At fixed “temperature”  $T$ , state occupation probability reaches Boltzman distribution

$$p(x) = \alpha e^{\frac{E(x)}{kT}}$$

$T$  decreased slowly enough  $\implies$  always reach best state  $x^*$   
because  $e^{\frac{E(x^*)}{kT}} / e^{\frac{E(x)}{kT}} = e^{\frac{E(x^*) - E(x)}{kT}} \gg 1$  for small  $T$

## Is this necessarily an interesting guarantee

Devised by Metropolis et al., 1953, for physical process modelling

Widely used in VLSI layout, airline scheduling, etc.

# Continuous state spaces

Suppose we want to site three airports in Romania:

- 6-D state space defined by  $(x_1, y_1), (x_2, y_2), (x_3, y_3)$
- objective function  $f(x_1, y_1, x_2, y_2, x_3, y_3) =$   
sum of squared distances from each city to nearest airport

**Discretization** methods turn continuous space into discrete space, e.g., **empirical gradient** considers  $\pm\delta$  change in each coordinate

**Gradient** methods compute

$$\nabla f = \left( \frac{\partial f}{\partial x_1}, \frac{\partial f}{\partial y_1}, \frac{\partial f}{\partial x_2}, \frac{\partial f}{\partial y_2}, \frac{\partial f}{\partial x_3}, \frac{\partial f}{\partial y_3} \right)$$

to increase/reduce  $f$ , e.g., by  $\mathbf{x} \leftarrow \mathbf{x} + \alpha \nabla f(\mathbf{x})$

Sometimes can solve for  $\nabla f(\mathbf{x}) = 0$  exactly (e.g., with one city).

**Newton–Raphson** (1664, 1690) iterates  $\mathbf{x} \leftarrow \mathbf{x} - \mathbf{H}_f^{-1}(\mathbf{x}) \nabla f(\mathbf{x})$   
to solve  $\nabla f(\mathbf{x}) = 0$ , where  $\mathbf{H}_{ij} = \partial^2 f / \partial x_i \partial x_j$